

Computer Aided Verification 2015

The SPIN model checker

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Material borrowed from Roberto Bruttomesso

1 Introduction

2 PROcess MEta LANGUAGE

- Data types
- Control structures
- Channels

1 Introduction

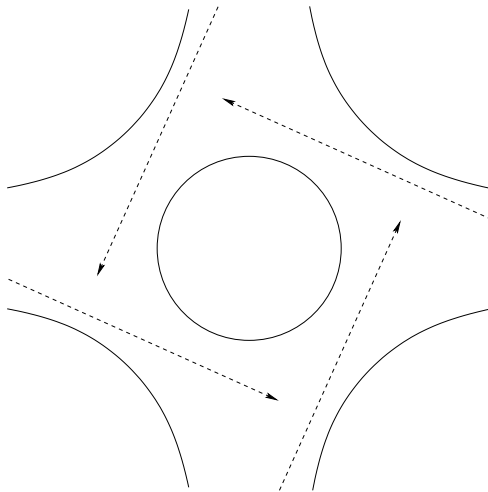
2 PROcess MEta LANGUAGE

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- **Explicit-state** model checker
 - works on-the-fly (no need to represent the entire structure of the system upfront)
- Targeted toward the verification of **concurrent** systems
 - real world situations involving multiple agents acting at the same time on the same resources
 - hardware or software with interacting components
- Input language: **PROMELA** (PRocess MEta LAnguage)
 - Designed for describing processes
 - Different from an imperative language
- Properties: **LTL** (Linear Temporal Logic)

Bugs in concurrent systems (1)

Circular blocking



Bugs in concurrent systems (2)

Deadly embrace

We have two global resources p (a printer) and s (a scanner), and two concurrent processes running in memory

process A

```
1  getPrinter( $p$ )  
2  getScanner( $s$ )  
3  ...  
4  releasePrinter()  
5  releaseScanner()  
end
```

process B

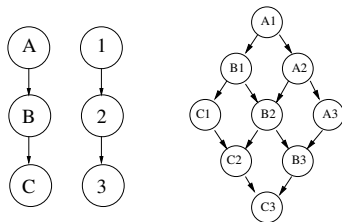
```
1  getScanner( $s$ )  
2  getPrinter( $p$ )  
3  ...  
4  releasePrinter()  
5  releaseScanner()  
end
```

The sequence A1, B1, A2, B2 generates a deadlock

Spin

Execution semantic

- In Spin we can specify processes, which are assumed to run **concurrently**
- Since we are dealing with multiple processes, at any point in time a number of instructions can be executed (unless deadlock)
- Spin chooses **non-deterministically** one instruction to be executed among the ones available



- Simulator (first half): runs a guided or a random simulation of the model defined by the user. Can be used to quickly check the behavior of a model
- Verifier (second half): generates a C program, optimized for performance, that **exhaustively** checks the validity of the property (or the absence of deadlocks)

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Data types

Basic types

Type	Range	Example
bit	$[0,1]$	<code>bit flag = 0;</code>
bool	false, true	<code>bool flag = false;</code>
byte	$[0,255]$	<code>byte b = 0;</code>
chan	$[1,255]$	see later ...
mtype	$[1,255]$	see later ...
pid	$[0,255]$	<code>pid p = 0;</code>
short	$[-2^{15}, 2^{15} - 1]$	<code>short s = 0;</code>
int	$[-2^{31}, 2^{31} - 1]$	<code>int i = 0;</code>
unsigned	$[0, 2^n - 1]$	<code>unsigned u : 8;</code>

Default value for uninitialized variables is “0”

Data types

Compound types

- Arrays (one dimension):

```
byte vec[10];
```

- Records:

```
typedef rec  
{  
    bit a;  
    byte b;  
}
```

- `mtype`: used to declare user-defined constants (similar to C `enum`)

Example:

```
mtype = { ack, nack, error };
```

Control structures

Processes

- In Promela we can only define processes
- `init` is a process which is automatically started at the beginning
- `proctype` can be used to define a process
- a process can be started with `run`
- `active` can be used to specify a `proctype` that is automatically started at the beginning

Control structures

Basic Statements

- Assignments: `count = count + 1`
 - **Always** executable: `count` is updated
- Conditions: `visitors == 10`
 - Executable **only if** they hold. Nothing is changed.
- Example: `busy.pml`
- Statements are separated by `;`
 - `x = x + 2; y = x + z`
- or by a `->` (usually when `condition -> assignment`)
 - `x == 0 ; y = x + z`
 - `x == 0 -> y = x + z`

Control structures

Case selection

```
if  
::  cond_1 -> list of statements 1  
::  cond_2 -> list of statements 2  
::  ...  
fi
```

- Only **one** option from the list will be executed
- It is chosen randomly among the satisfied conditions
- If no condition holds, the process blocks until a condition becomes **true**
- Example `if.pml`

Control structures

Repetition

```
do
::  cond_1 -> list of statements 1
::  cond_2 -> list of statements 2
::  ...
od
```

- Similar to `if`, but repeated over time
- Only **one** option from the list will be executed
- It is chosen randomly among the satisfied conditions
- If no condition holds, the process blocks until a condition becomes `true`
- Example `do.pml`

Control structures

Other useful operators

- **else**: in **if-fi** or **do-od** statements become true if any other condition fails
- **goto ...**: can be used to jump to another process location
- **break**: force exit from a **do-od** loop
- **timeout**: becomes true if no other statement can be executed
- **atomic { ... }** : process a list of statements without interleaving
- **skip**: does nothing

Example

```
int iter;

proctype counter( byte count )
{
    do
        :: (count != 0) ->
            if
                :: count = count + 1
                :: count = count - 1
            fi
        :: (count == 0) -> break
    od;
    printf( "Iterations done:  %d" iter );
}
```

Exercise 1

Collatz Conjecture ($3x+1$ Problem)

- Let n be a natural number. Consider the following loop:

1	while $n \neq 1$
2	if n is even then $n := n/2$
3	else $n := 3 \cdot n + 1$

- Example: $n = 5$ generates the sequence, 5, 16, 8, 4, 2, 1
- Exercise:

Implement the loop in PROMELA and return as output the number of iterations done

(use “%” to compute the remainder)

Exercise 1

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- Exercise:

Implement the loop in PROMELA and return as output the number of iterations done

extra Write a routine that computes the number between 1 and 100 that takes the highest number of steps

(use “%” to compute the remainder)

Communication

Channels (1)

- Data can be exchanged via global variables (as seen so far) or using **channels**
- a channel is FIFO queue

Example declaration:

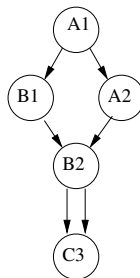
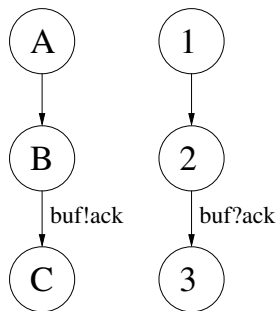
```
chan c = [2] of { byte }
```

Operations:

- **send** an element **e** into the channel **c**
 - Syntax: **c!e**
 - Executable if the channel is not full
- **receive** an element from the channel **c**, and store it into **e**
 - Syntax: **c?e**
 - Executable if the channel is not empty
- Example `prod_cons.pml`

- **Empty** channels can be declared to define **rendez-vous** points

```
chan buf = [0] of { mtype };  
mtype { ack, nack };
```



Exercise 2

Modify `prod_cons.pml` as follows:

- `producer()` sends odd and even numbers into the buffer
- write `consumer_odd()` that prints only odd numbers
- write `consumer_even()` that prints only even numbers
- `producer()` should send termination signal to both consumers
- every number must be printed