

Basics of Complexity Analysis: The RAM Model and the Growth of Functions

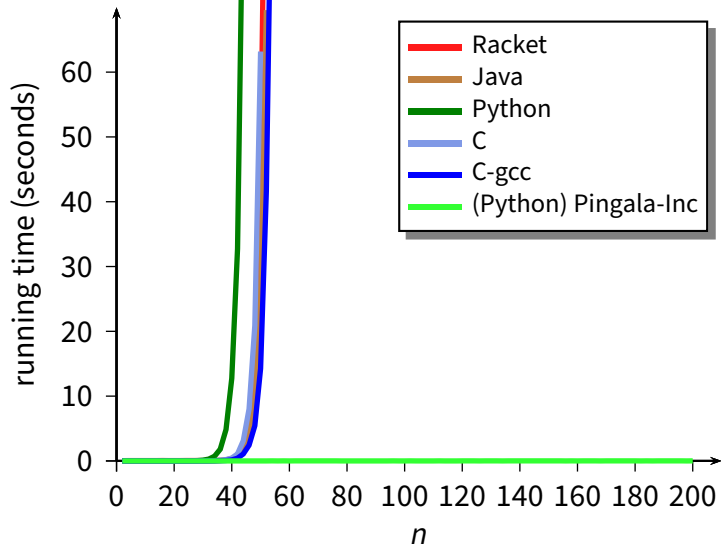
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February 22, 2024

- Informal analysis of two Pingala algorithms
- The *random-access machine* model
- Measure of complexity
- Characterizing functions with their asymptotic behavior
- Big-O, omega, and theta notations

Slow vs. Fast Pingala



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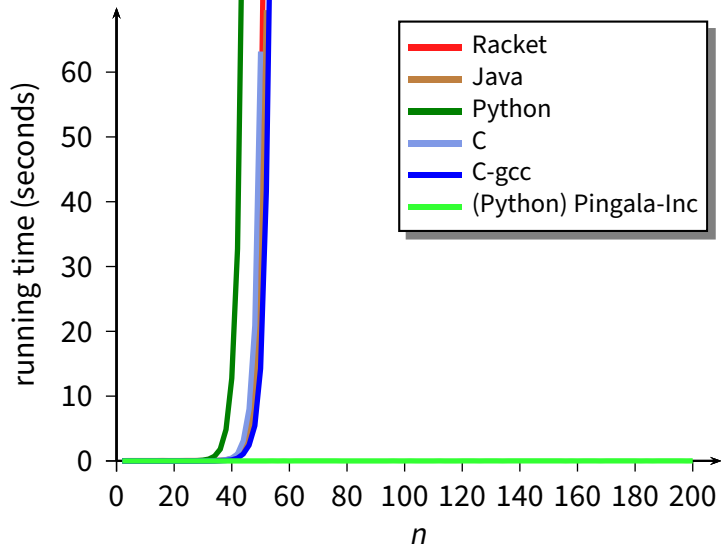
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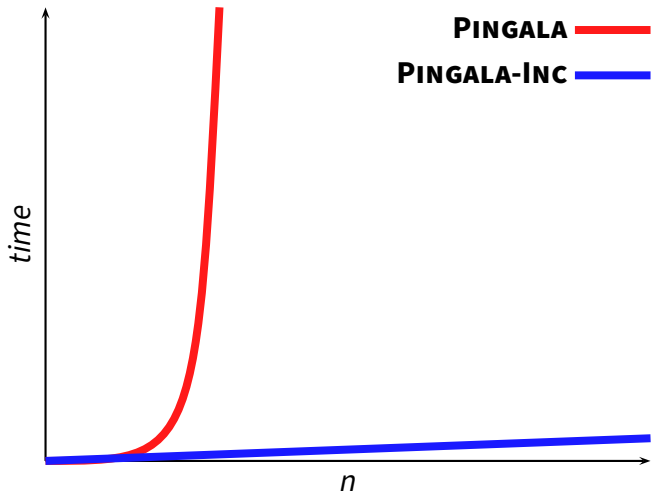
▶ in a way that is *specific to the algorithms*

▶ but *independent of implementation details*

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- A ***basic step*** in the RAM model takes a ***constant time***

Analysis in the RAM Model

PINGALA-INC(n)

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1  if  $n \leq 2$ 
2      return  $n$ 
3   $pprev = 1$ 
4   $prev = 2$ 
5  for  $i = 3$  to  $n$ 
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c_1	1
c_2	0
c_3	1
c_4	1
c_5	$n - 1$
c_6	$n - 2$
c_7	$n - 2$
c_8	$n - 2$
c_9	1

$$T(n) = c_1 + c_3 + c_4 + c_9 + (n - 1)c_5 + (n - 2)(c_6 + c_7 + c_8)$$

- Does a load/store operation cost more than, say, an arithmetic operation?

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- ***We do not care about the specific costs of each basic step***
 - ▶ these costs are likely to vary significantly with languages, implementations, and processors
 - ▶ we simplify our model by effectively considering only the maximal cost of any basic step
 - ▶ so, we assume $c_1 = c_2 = c_3 = \dots = c_i$

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$$T(n) = nC_1 + C_2 \Rightarrow T(n) \text{ is a linear function of } n$$

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- A ***basic step*** in the RAM model takes a ***constant time***

- ▶ “constant” means ***independent of the input size***

- The *specific* constant is a ***technological factor***

- ***Technology changes***

- ...so ***we ignore any specific multiplicative or additive constants***

- ...effectively we allow for ***any scaling factor***

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$$T(n) = Cn$$

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FINDEQUALS(A)

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$$T(n) = C \frac{n(n-1)}{2}$$

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 - ▶ so we ignore lower-order terms

Example:

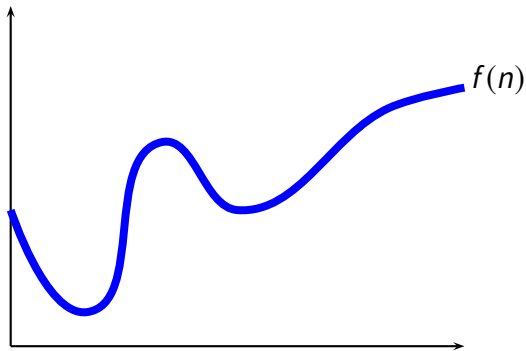
Algorithm 1 costs $T_1(n) = 100n + 3000$ basic steps

Algorithm 2 costs $T_2(n) = 0.02n^2 + 2$ basic steps

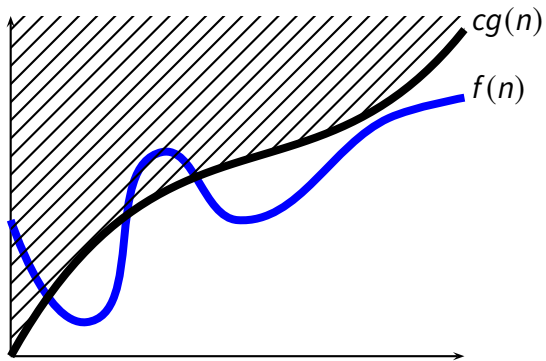
Which is best?

- Given a function $g(n)$, we define the *family of functions* $O(g(n))$

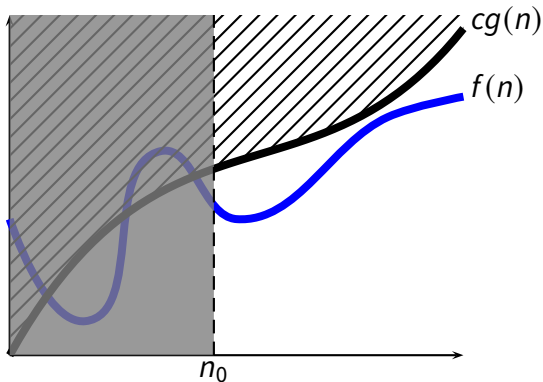
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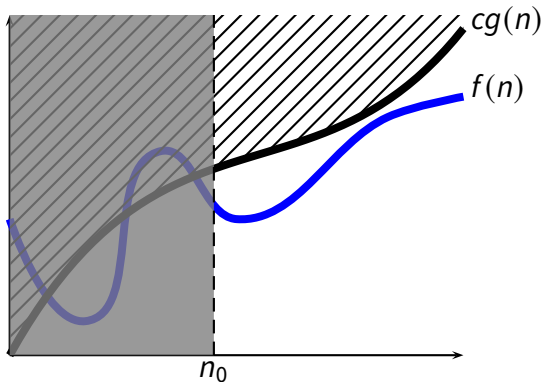


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$$O(g(n)) = \{f(n) : \exists c > 0, \exists n_0 > 0 \\ : 0 \leq f(n) \leq cg(n) \text{ for all } n \geq n_0\}$$

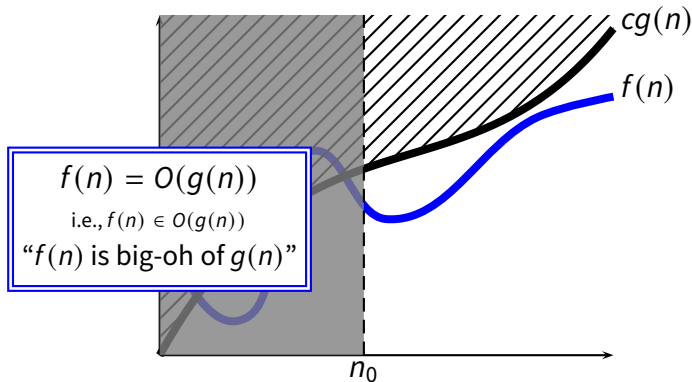
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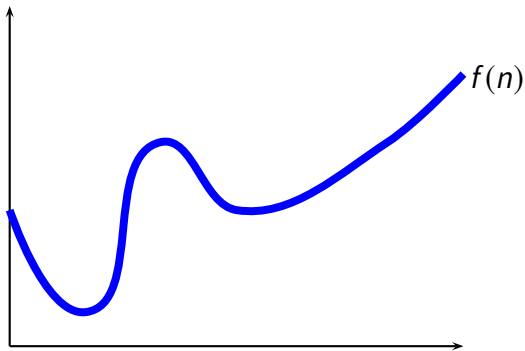
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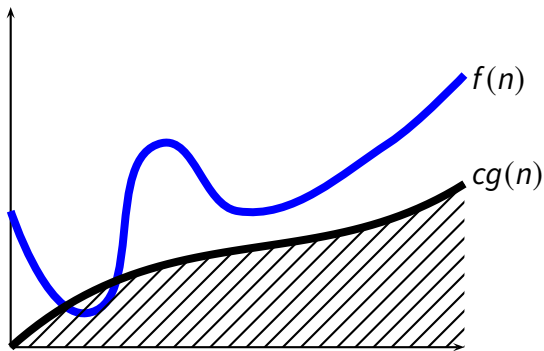
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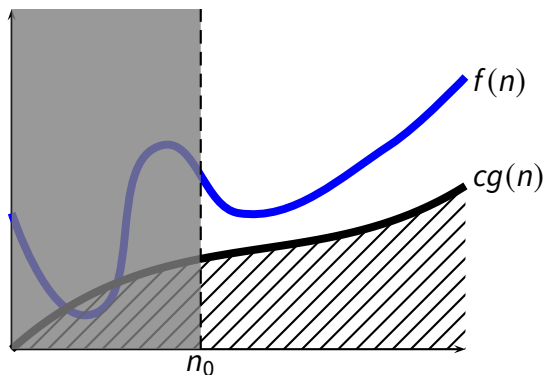
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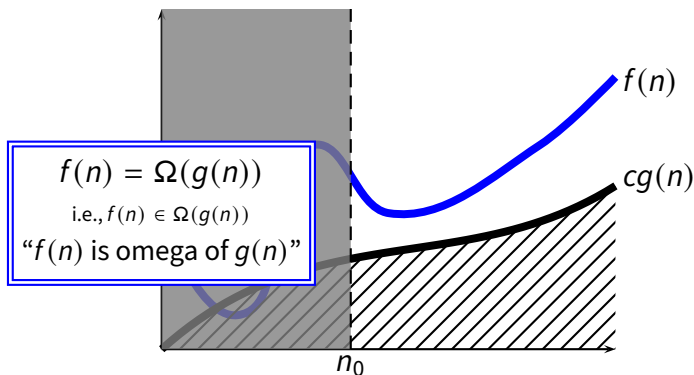


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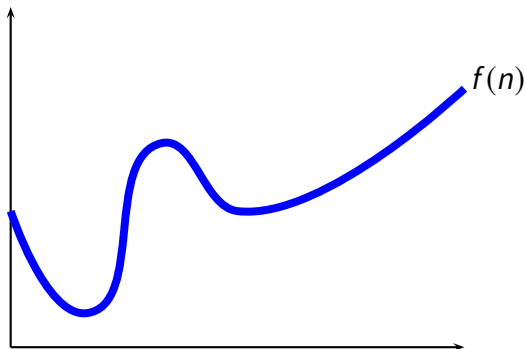
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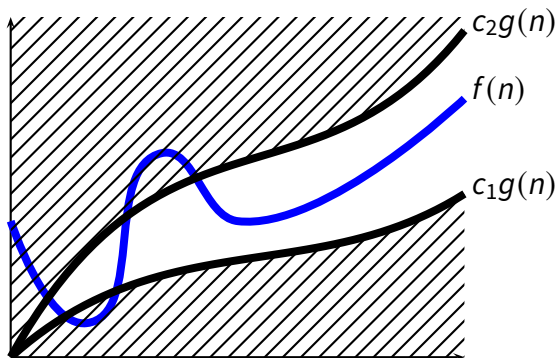
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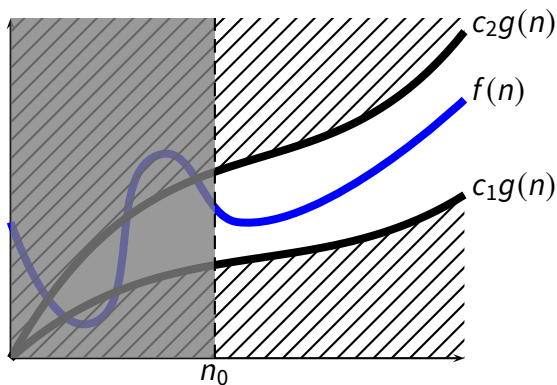
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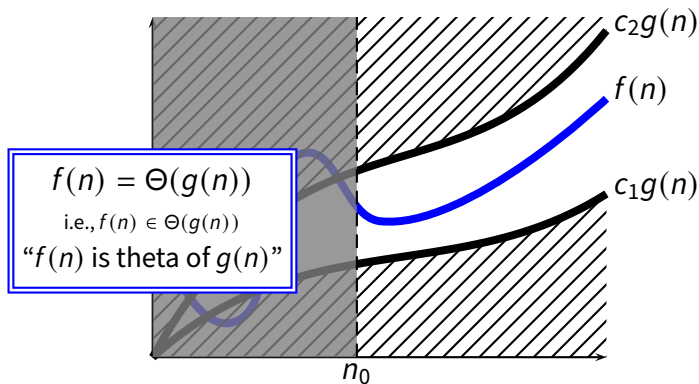


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In fact, the fundamental *prime number theorem* says that

$$\lim_{n \rightarrow \infty} \frac{\pi(n) \ln n}{n} = 1$$

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- $T(n) =$ complexity of **PINGALA-INC**

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- $T(n) = \text{complexity of PINGALA-INC} \Rightarrow T(n) = \Theta(n)$
- We characterize the behavior of $T(n)$ *in the limit*
- The Θ -notation is an **asymptotic notation**

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- So, what is the complexity of **FINDEQUALS**?

FINDEQUALS(*A*)

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1  for i = 1 to length(A) - 1
2      for j = i + 1 to length(A)
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$$T(n) = \Theta(n^2)$$

- ▶ $n = \text{length}(A)$ is the *size of the input*
- ▶ we measure the *worst-case complexity*

Θ , O , and Ω as Relations

- *Theorem:* for any two functions $f(n)$ and $g(n)$,
 $f(n) = \Omega(g(n)) \wedge f(n) = O(g(n)) \Leftrightarrow f(n) = \Theta(g(n))$

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Θ , O , and Ω as Anonymous Functions

- We can use the Θ -, O -, and Ω -notation to represent anonymous (unknown or unspecified) functions
E.g.,

$$f(n) = 10n^2 + O(n)$$

means that $f(n)$ is equal to $10n^2$ plus a function we don't know or we don't care to know that is asymptotically at most linear in n .

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