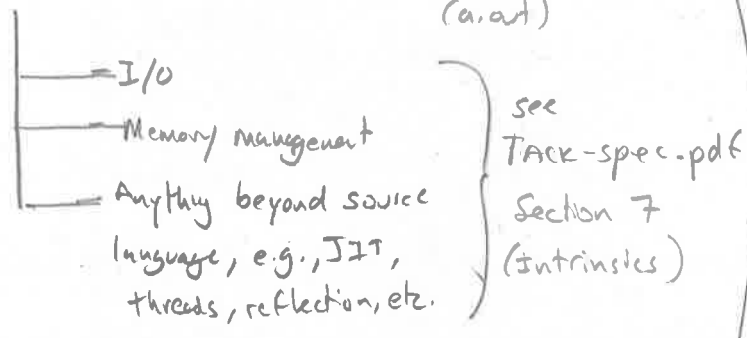
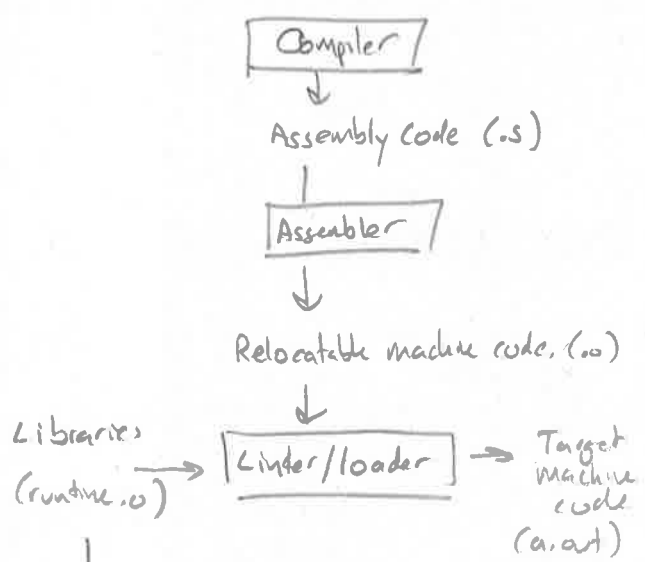


Runtime environments pg. 1

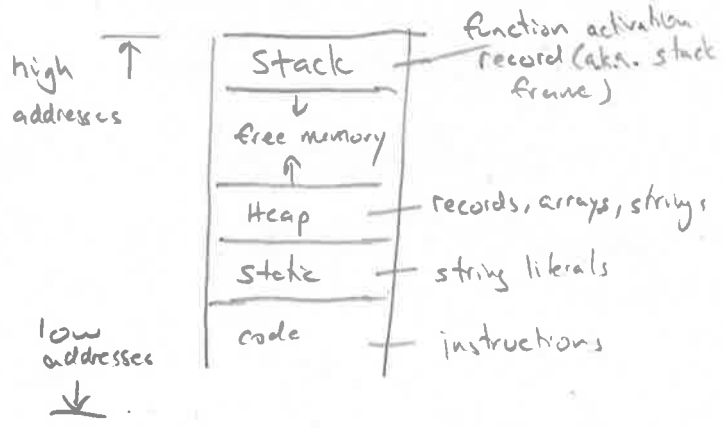
Lecture topics

- I. Runtime Systems
- II Stack
- III Heap

I. Runtime Systems



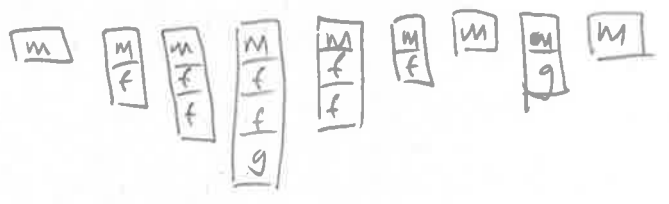
Storage organization



II Stack

```

main { f(); g(); }
f { if (... f()); else g(); }
g { ... }
  
```



Example TACK

$$x = f(E_0, \dots, E_{n-1})$$

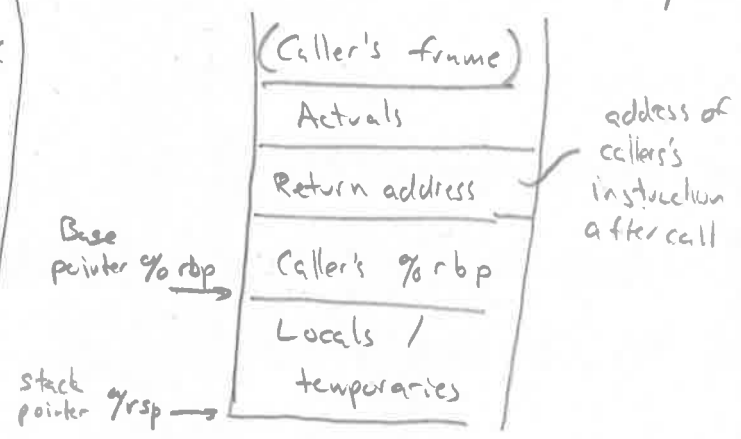
IR

E_0 . code
 ...
 E_{n-1} . code
 param[0:n] = E_0 .addr
 ...
 param[n-1:n] = E_{n-1} .addr
 x = call f on

→ E;

E.code
return E.addr

Stack frame layout (see x64-intro.pdf)



Runtime environments pg. 2

Call sequence

- put actuals in register/stack
 - call (pushes return address)
- } caller
- push %rbp
 - move %rbp, %rsp
 - sub %rsp, frame size
- } callee

Return sequence

- move %rax, return value
 - move %rsp, %rbp
 - pop %rbp
 - ret (pops return address)
- } callee
- remove actuals from stack
 - use %rax
- } caller

Other stack topics:

- saving registers across calls
- variable size data on stack
- accessing non-local stack variables

III Heap

x = (a=1);

y = (a=2);

x = null

→ y;



lifetimes of heap objects are not nested, may be indefinite

	Java	C	C++	Tack
Allocation	new	malloc	new	records/array literals
Deallocation	garbage collection	free	delete	unspecified