The Network Layer

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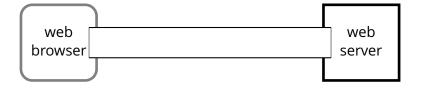
December 1, 2016

Outline

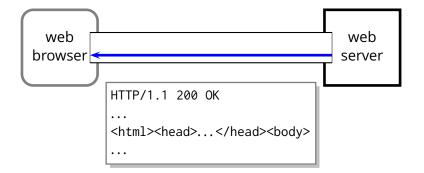
- Basic network-layer architecture of a datagram network
- Introduction to forwarding
- Introduction to routing
- General architecture of a router
- Switching fabric and queuing
- Internet network-layer protocol
- The Internet protocol (IP)
 - Fragmentation



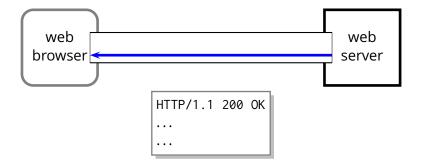
web server



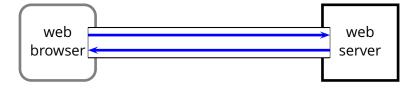








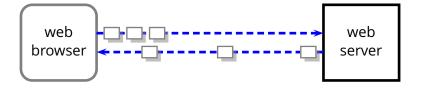
Transport Level



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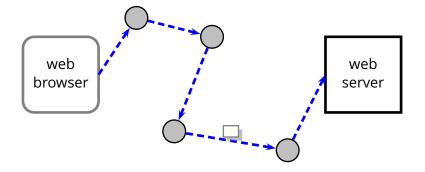


Network Layer

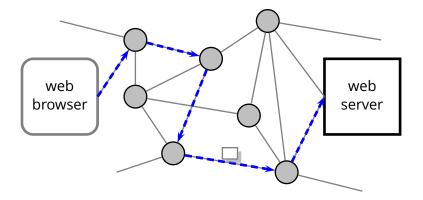


web server

Network Layer



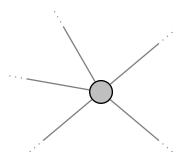
Network Layer



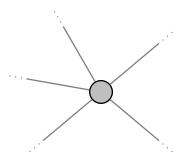






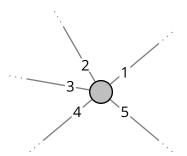


Fundamental component of the network layer



■ Fundamental component of the network layer

A node in a graph



- Fundamental component of the network layer
- A node in a graph
- A finite set of input/output (physical) connections
 - a.k.a., *interfaces* or *ports*

Packet-switched network

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 - information is transmitted in discrete units called *datagrams*

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"Best-effort" service

delivery guarantee: none

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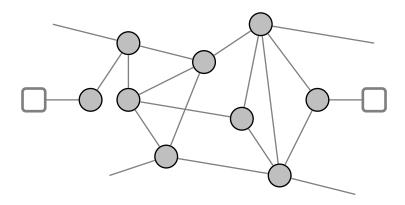
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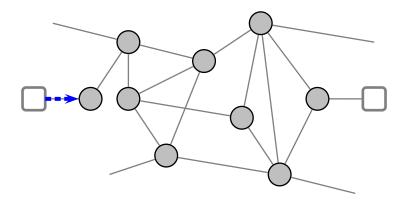
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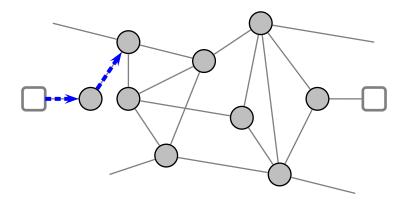
Connectionless service

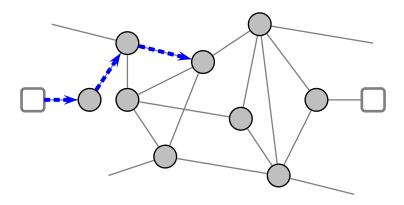
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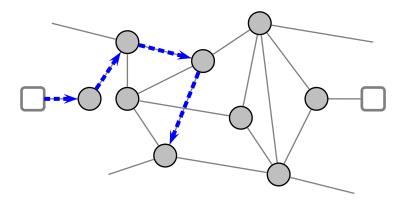
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- congestion indication: none

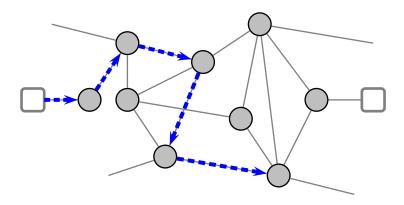


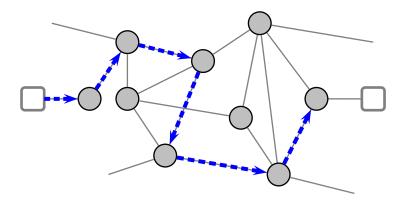


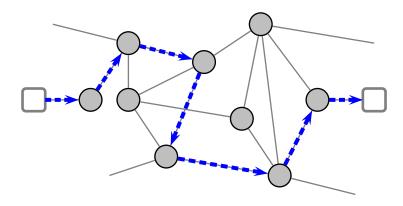


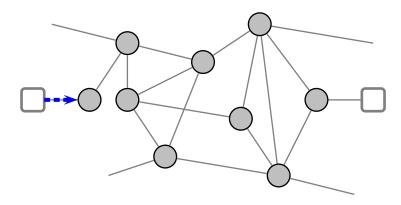


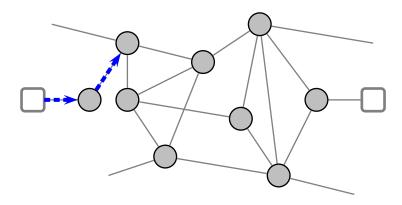


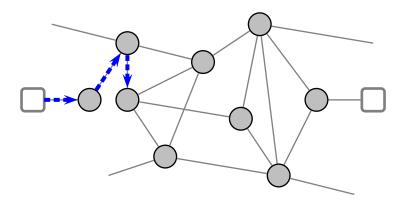


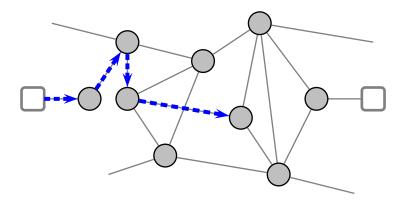


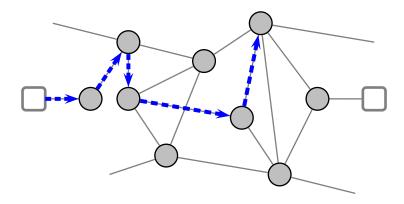


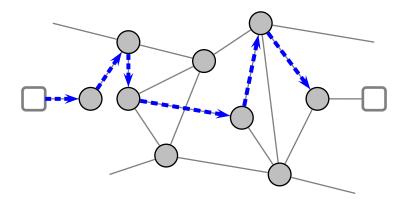


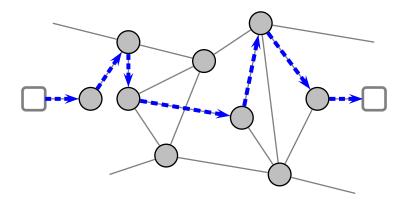


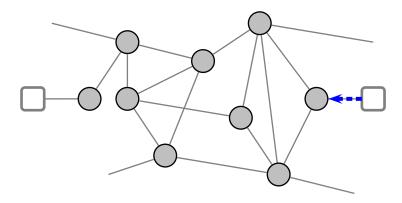




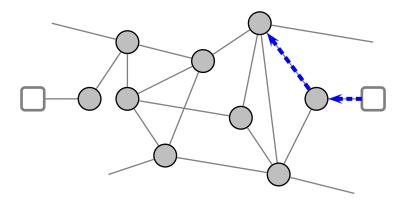




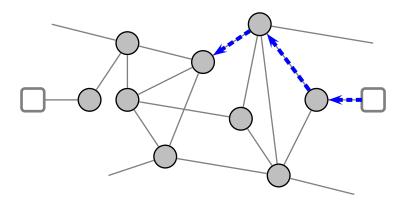




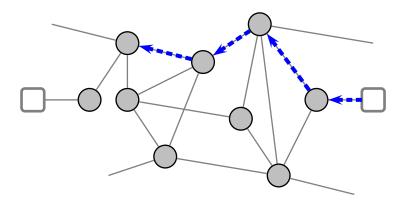
Potentially *multiple paths* for the same source/destination



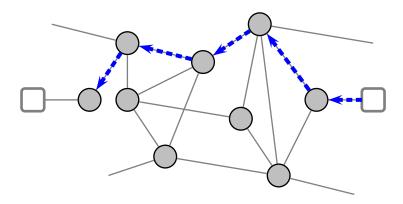
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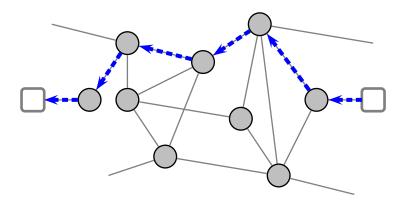
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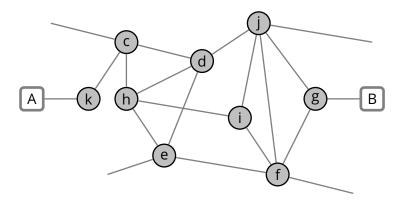
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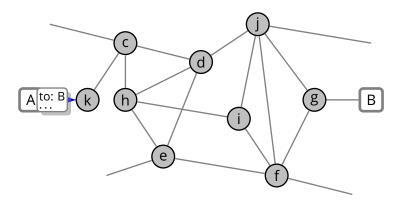
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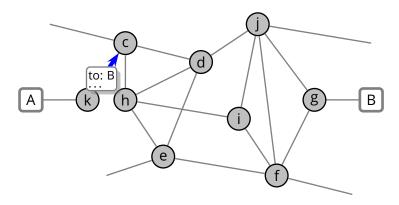


A sends a datagram to B



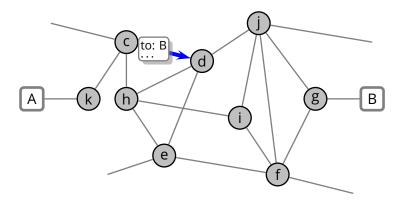
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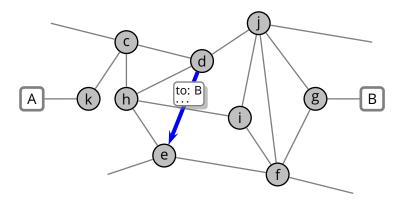
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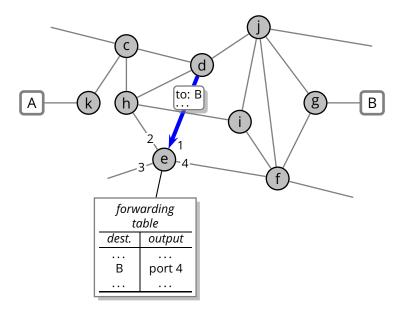


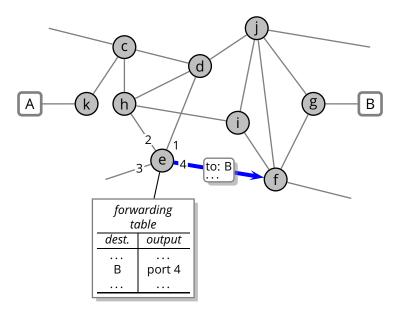
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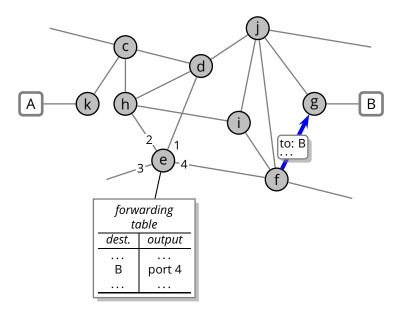
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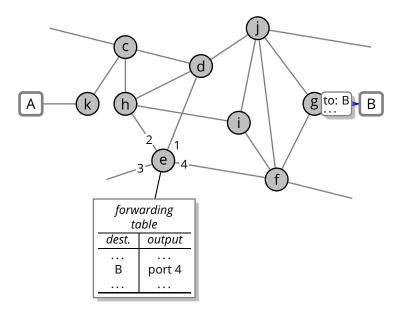












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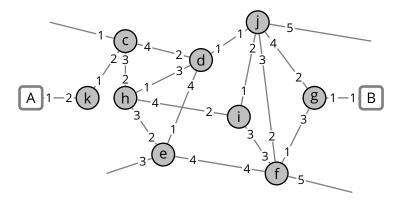
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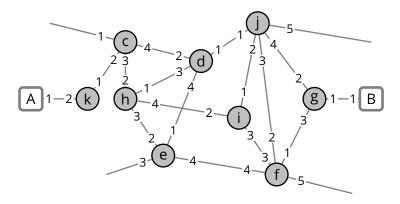
- how big is the forwarding table?
- how fast does the router have to forward datagrams?
- how does the router build and maintain the forwarding table?

Routing

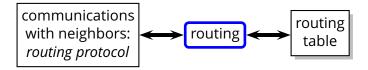
Routing

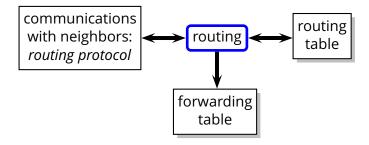


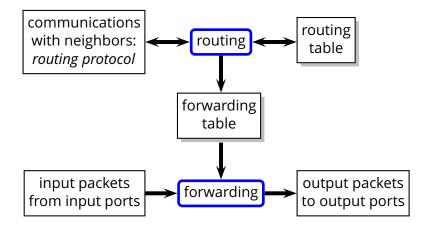
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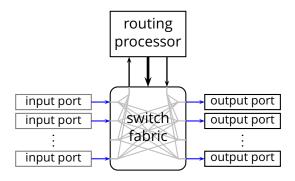


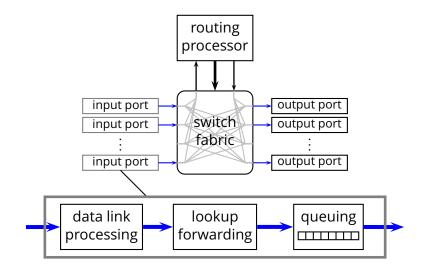
rοι	router k				
А	2				
В	1				

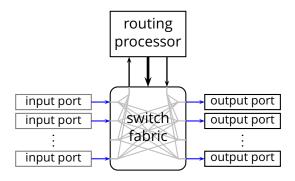


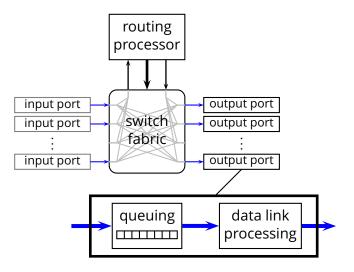














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- Input ports
 - queuing may occur here if the switching fabric is slower than the aggregate speed of all the input lines. I.e., $R_S < nR_{in}$
- Output ports
 - queuing may occur here because of the limited throughput of the output link. I.e., R_{out} < min(R_s, nR_{in})

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- Deciding when to drop packets, and which packets to drop
 - *drop tail:* drop arriving packets when queues are full
 - active queue management: a set of policies and algorithms to decide when and how to drop or mark packets in the attempt to prevent congestion

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- OSPF
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- datagram format
- fragmentation and packet handling

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ICMP

- error reporting
- signaling

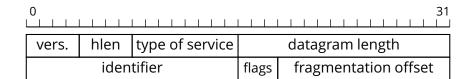


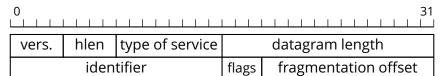








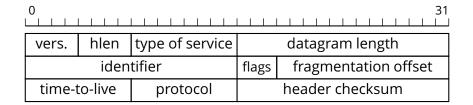




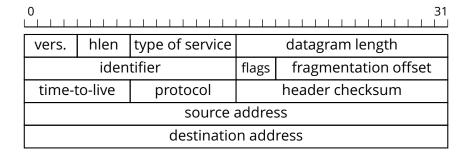
time-to-live



vers.	hlen	type of service		datagram length
identifier		flags	fragmentation offset	
time-t	o-live	protocol		



IPv4 Datagram Format

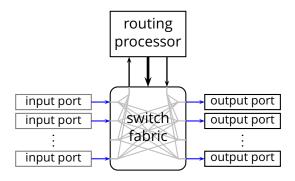


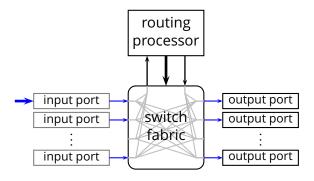
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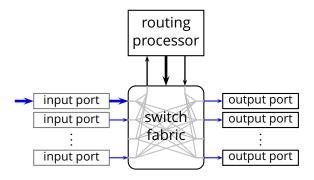
0				31		
vers.	hlen	type of service	datagram length			
identifier			flags fragmentation offset			
time-to-live protocol header checksum			header checksum			
		source	addre	SS		
		destinatio	n add	ress		
options (if any)						

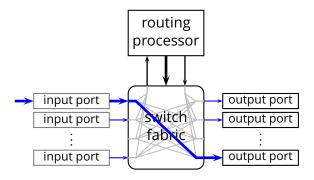
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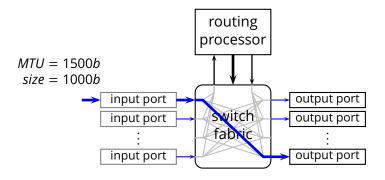
0				31		
vers.	hlen	type of service		datagram length		
identifier flags fragmentation offset						
time-to-live protocol header checksum						
source address						
	destination address					
options (if any)						
	data					

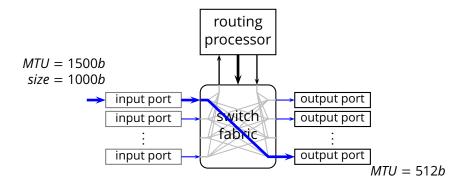


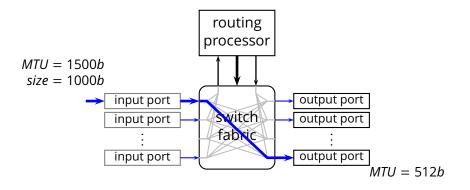




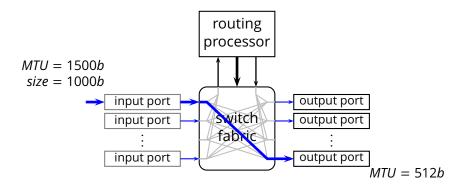








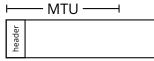
How does the router handle cases where the size of an input datagram exceeds the maximum transmission unit (MTU) of the output link?



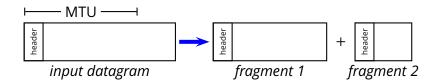
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- The datagram is *fragmented*

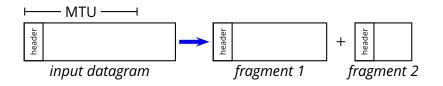
header					
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input datagram

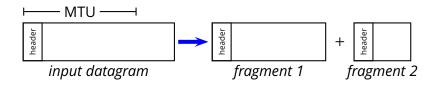


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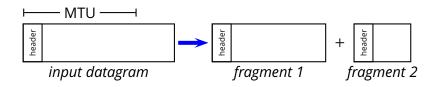


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Requirements

- destination must recognize two fragments of the same original datagram
- destination must see if and when all the fragments have been received
- intermediate routers must be able to fragment a datagram to whatever level necessary

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identifier	fragment	more	header	total
	offset	fragments	length	length
789	0	0	20	1020

Fragmentation to an MTU of 512

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identifier	fragment	more	header	total
	offset	fragments	length	length
789	0	1	20	508

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identifier	fragment	more	header	total
	offset	fragments	length	length
789	61	1	20	508

Fragmentation to an MTU of 512

identifier	fragment	more	header	total
	offset	fragments	length	length
789	0	1	20	508

identifier	fragment	more	header	total
	offset	fragments	length	length
789	61	1	20	508

identifier	fragment	more	header	total
	offset	fragments	length	length
789	122	0	20	44